



# Firewall Presentation

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# Disclaimer



This presentation:

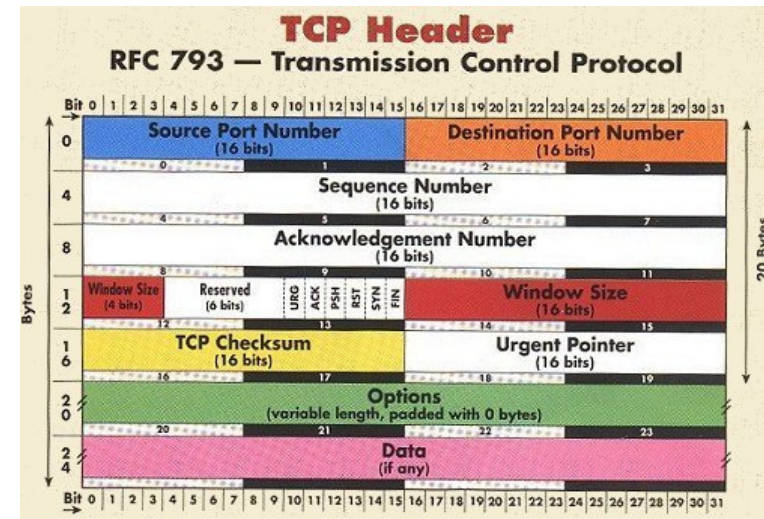
- does not contain NRC official positions
- is not guidance on how to configure firewalls
- is an overview of firewalls and their limitations
- is a demonstration of how attackers can bypass firewalls

# What is a Firewall?

Q: What is a firewall?

A: A firewall is a computer.

- A firewall has the following:
  - Two or more network cards
  - Processor, RAM, hard drive
  - Operating System



Q: What makes a firewall different from other computers?

A: Very little.

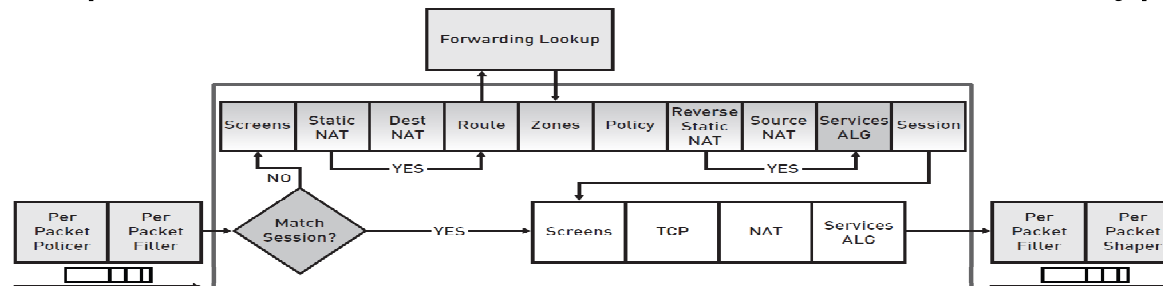
- Designed to analyze and filter data flows at its most basic level
- May include additional logic to perform real-time contextual analysis of data flows
- May include specialized networking hardware to aid in this task

# What is a Firewall?

Q: What is the purpose of a firewall?

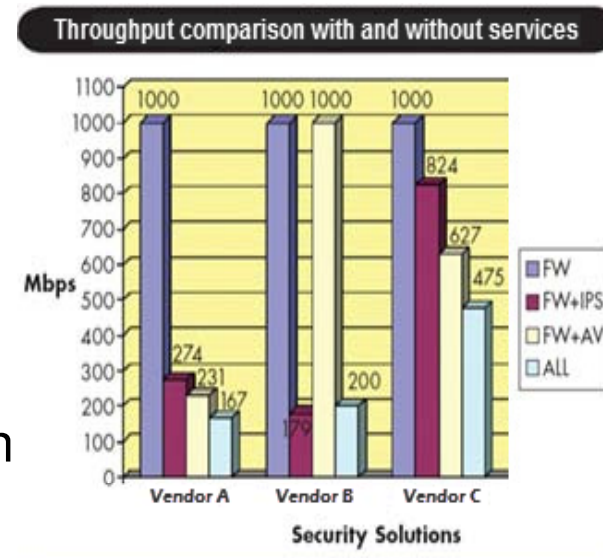
A: To control the flow of data between networks according to pre-defined rules

- Packet Filtering (by port, by protocol, by source address, by destination address)
- Stateful Inspection (can determine if a packet is part of an existing data flow)
- Other features include the following:
  - “Application Aware:” contains logic specific to common application (web, FTP, Secure Shell, etc.)
  - Quality of Service: Traffic prioritization and scheduling
  - Session Inspection: Can search a data flow for certain types of content



# Firewall Limitations

- A firewall cannot perform all security tasks
  - Hardware limitations
  - Memory and overhead limitations
  - Time limitations
  - Logic limitations
  - Encrypted traffic payloads are not visible
  - Firewalls do not typically do traffic normalization



- As a computer, a firewall can have vulnerabilities
  - CVE-2012-4661: Multiple Vulnerabilities in Cisco ASA 5500 Series Adaptive Security Appliances and Cisco Catalyst 6500 Series ASA Services Module
  - CVE-2012-5316: Multiple cross-site scripting (XSS) vulnerabilities in Barracuda Spam & Virus Firewall 600
  - **ICSA-12-102-05: Siemens Scalance S Multiple Security Vulnerabilities**

# Firewall Limitations

A firewall is only as good as its ruleset.

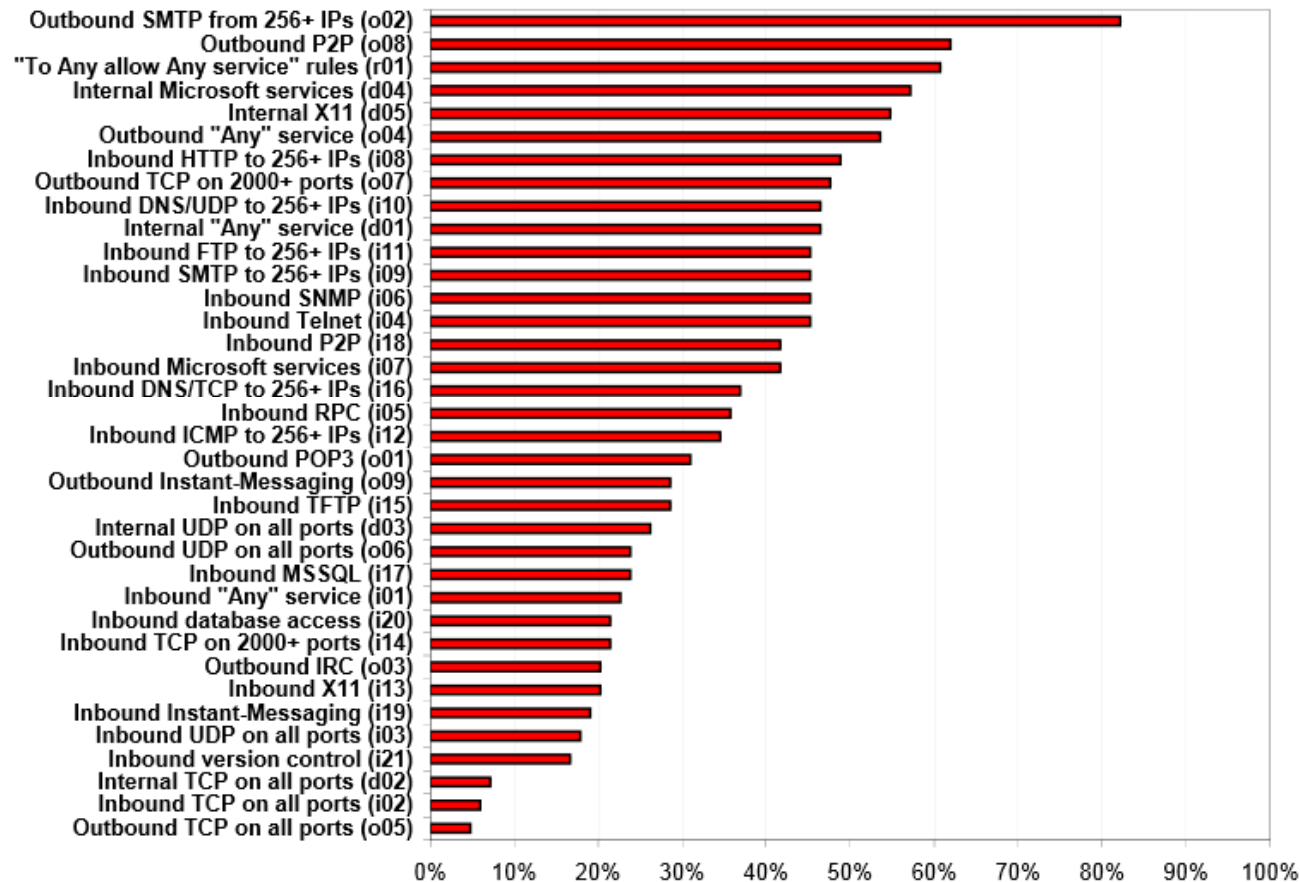
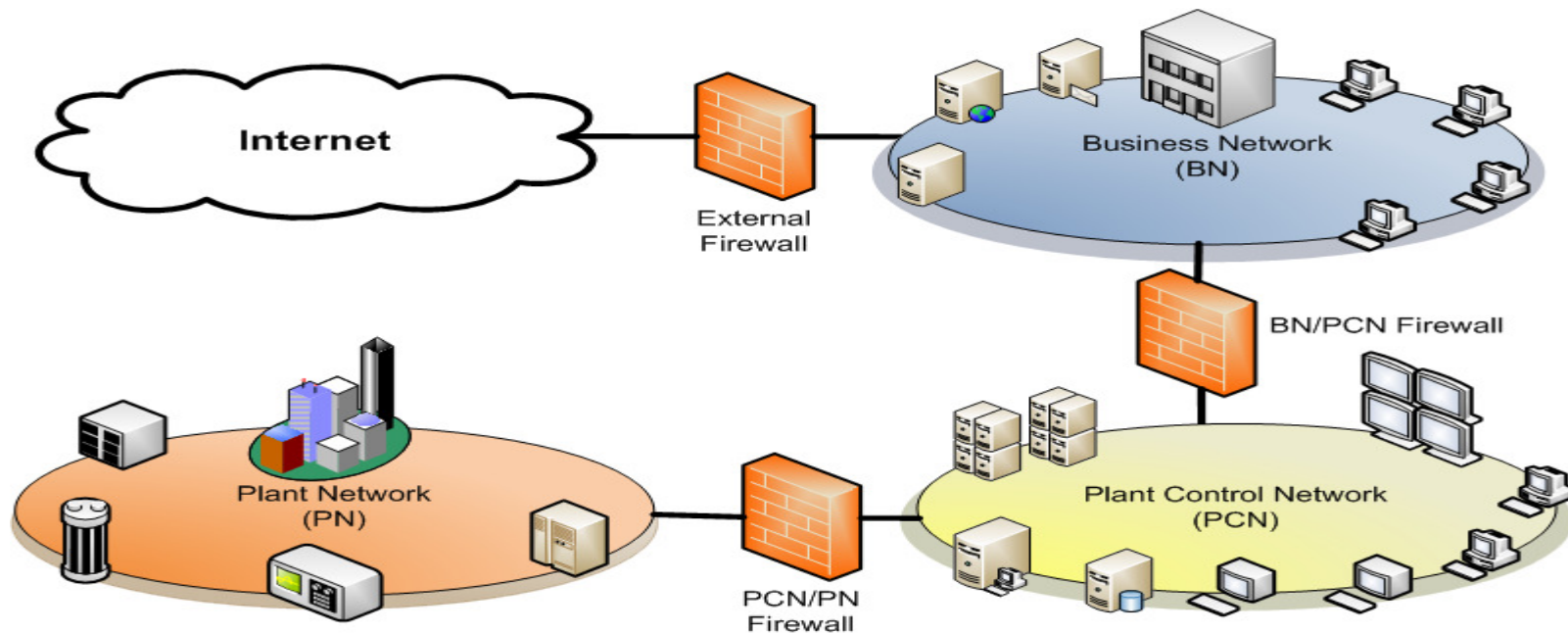


Figure 2: Distribution of configuration errors.

# Typical Network Architecture

- Business network acts as backbone
- Firewall between business network (BN) and plant control network (PCN)
- Firewall between PCN and plant network (PN) may or may not be in place

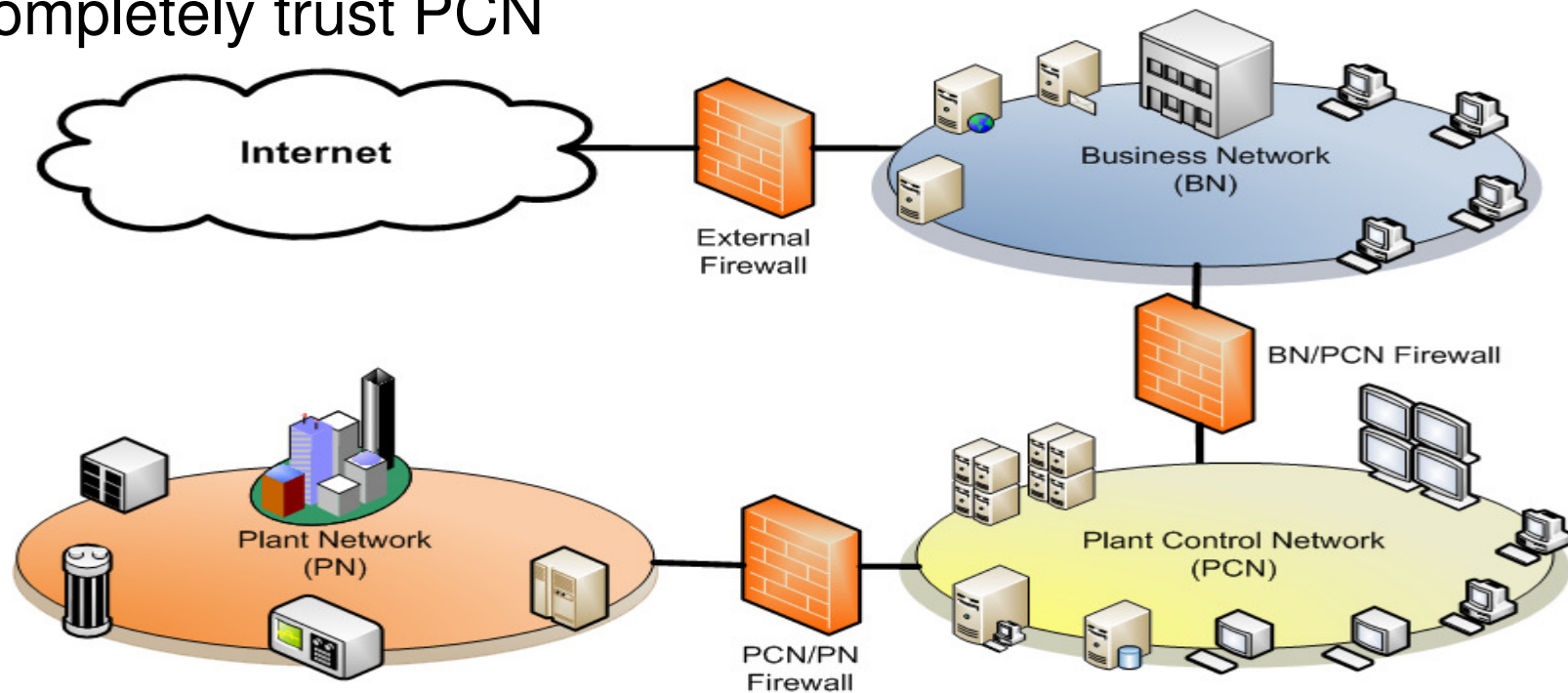




# Typical Network Architecture

## Problems:

- BN/PCN Firewall is configured to partially or completely trust BN
- PCN/PN Firewall is configured to partially or completely trust PCN





# Common Weaknesses to Model



- Poorly configured firewalls (historical, political, or legacy technical reasons)
  - Passing Microsoft Windows networking packets
  - Passing remote services (rsh, rlogin)
  - PCN/PN having trusted hosts on the business LAN
  - Not providing outbound data rules
- Peer links that bypass or route through external firewall direct to PCN or PN

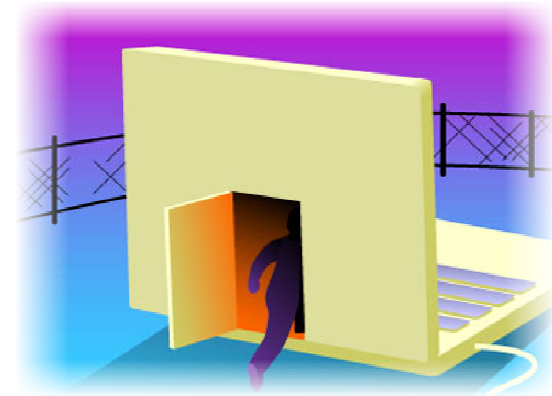
# Common Weaknesses to Model



- IT controlled assets in the PCN or PN (communications links, replicated services)
- Vendor links for remote maintenance/monitoring
- Out-of-band communications channels (backup links to RTUs)

# Getting Inside the Trusted Network

- Passive Evasion - The victim “phones home” to the attacker
  1. Phishing/spearphishing
  2. Malicious website/drive-by infection
  3. “Sneakernet” infection
  4. Social Engineering
- Indirect Evasion – Traffic appears to be authentic
  1. Stolen remote access credentials
  2. VPN piggyback
  3. Session hijacking
  4. Address spoofing (for internal zones)

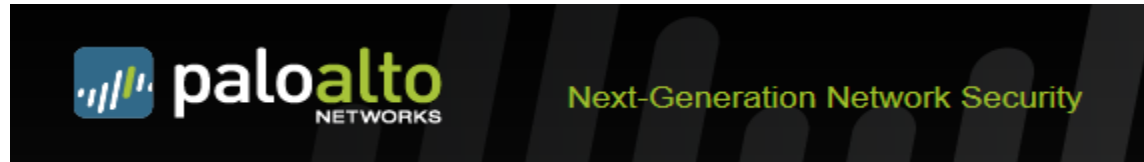


# Getting Inside the Trusted Network



- Active Evasion
  1. Attack exposed services (Web, E-mail)
  2. Attack firewall vulnerabilities
  3. Exploit weak ruleset/poor configuration
  4. “Trick” or subvert the firewall logic with protocol manipulation (AET)
  5. Find out-of-band channels (wireless, modems, satellite links)
  6. Get physical access to firewall or other infrastructure

# Case Study – Palo Alto Networks



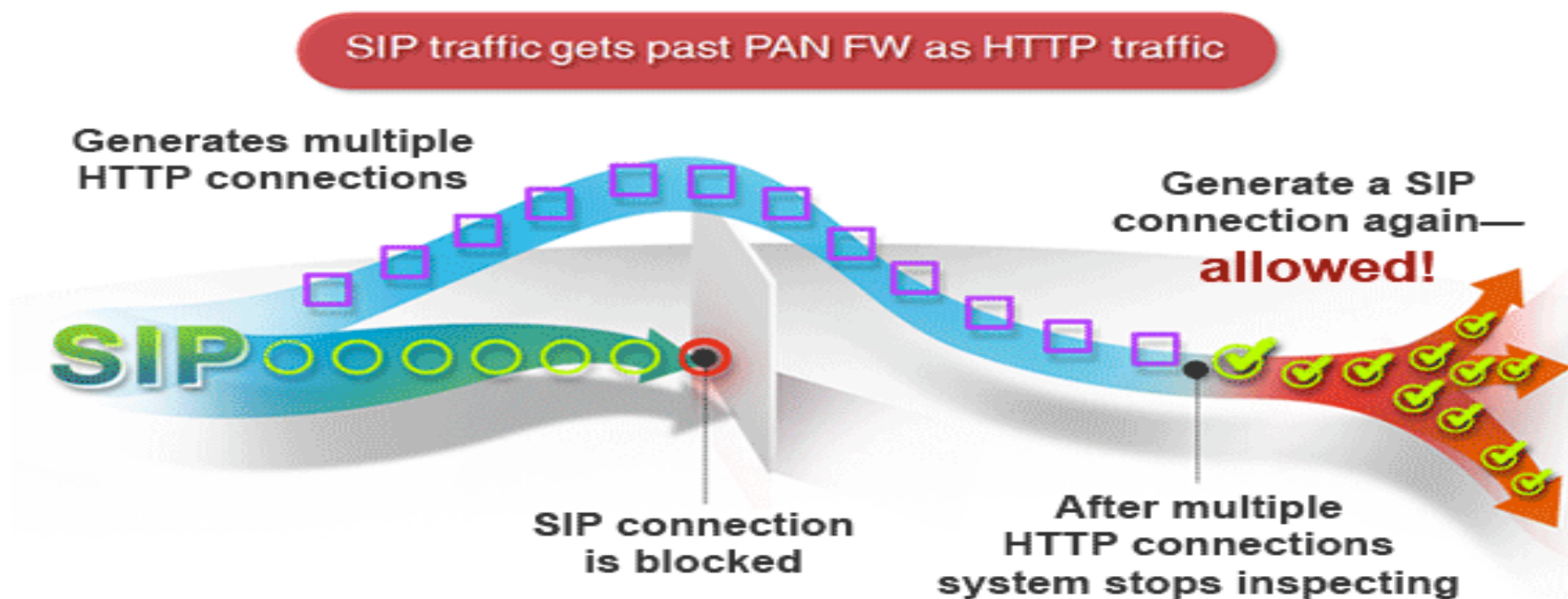
- Founded in 2005 by Checkpoint veteran
- First firewall product developed in 2007
- First of the “Next Generation” firewalls<sup>1</sup>
- Named leader in the 2011 Gartner “Magic Quadrant” report<sup>2</sup>
- At Defcon 19 (Dec 2011), Palo Alto firewall demonstrated to have fatal design flaw

1. Pescatore, J. & Young, G. (2009, October 19). Defining the Next-Generation Firewall. Gartner RAS Core Research Group. Retrieved from: <http://img1.custompublish.com/getfile.php/1434855.1861.sqqycbrdwq/Defining+the+Next-Generation+Firewall.pdf>, retrieved 2012-12-02
2. Denne, S. (2011, December 16). Palo Alto Networks hits the Magic Quadrant for firewalls. The Wall Street Journal. Retrieved from: <http://blogs.wsj.com/venturecapital/2011/12/16/palo-alto-networks-hits-the-magic-quadrant-for-firewalls/>
3. Woodberg, B. (2011). Palo Alto Networks Security Bypass. Defcon 19. Retrieved from: <http://www.youtube.com/watch?v=AuaCrRIIqnQ>

# Case Study – Palo Alto Networks

Cache poisoning attack:

- HTTP port open, SIP port blocked
- Attacker generates large number of HTTP sessions
- Memory cache fills, traffic no longer inspected
- HTTP session re-established as SIP, bypassing filter

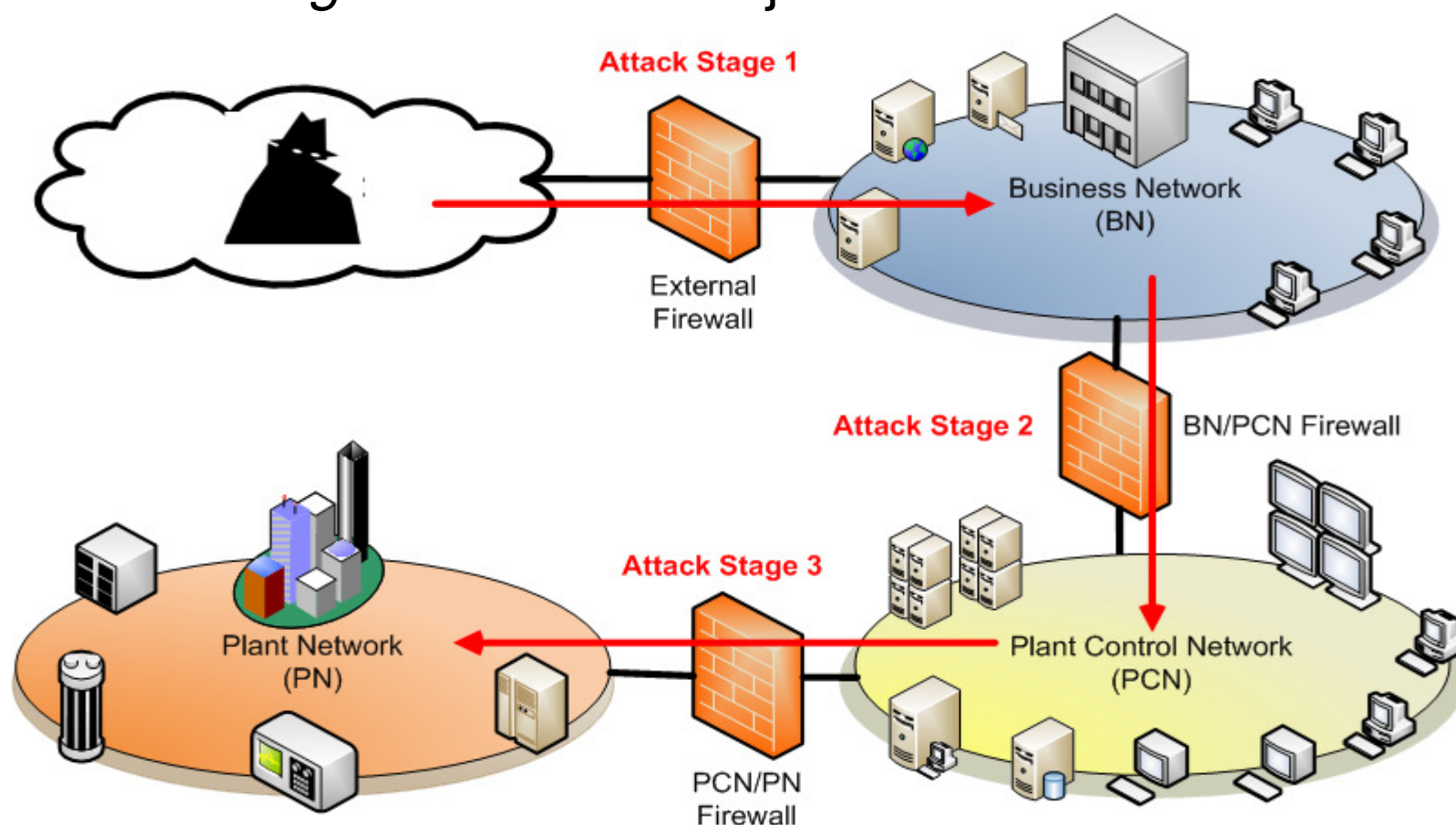


# Demonstration

*Attack Stage 1* – Desktop attack

*Attack Stage 2* – Impersonation Attack

*Attack Stage 3* – Session Hijack





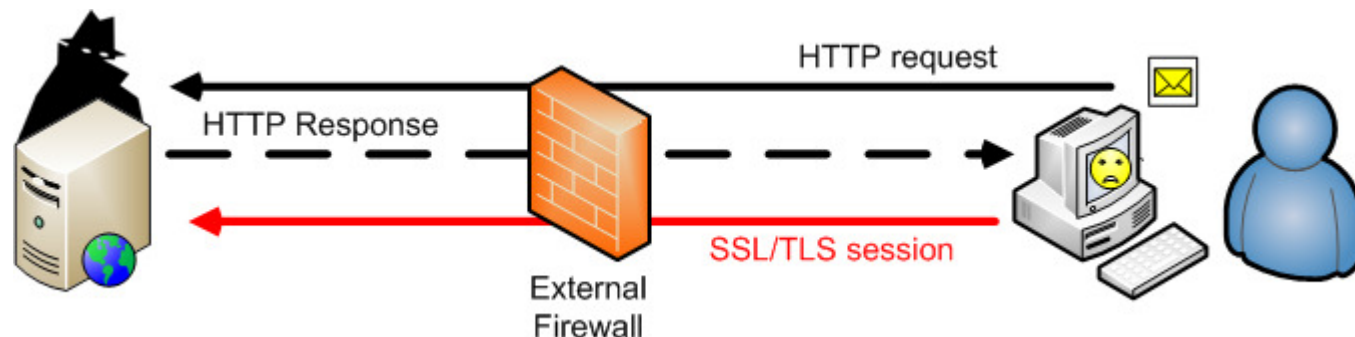
# Attack Stage 1– Desktop Attack

## Scenario 1:

- Attacker crafts email message to employee
  - Looks very believable, may come from spoofed address of trusted source
- Email contains link to compromised website

## Scenario 2:

- Employee goes to trusted website, which has link to infected website, employees computer is infected without knowledge (watering hole attack)



# Attack Stage 1– Desktop Attack

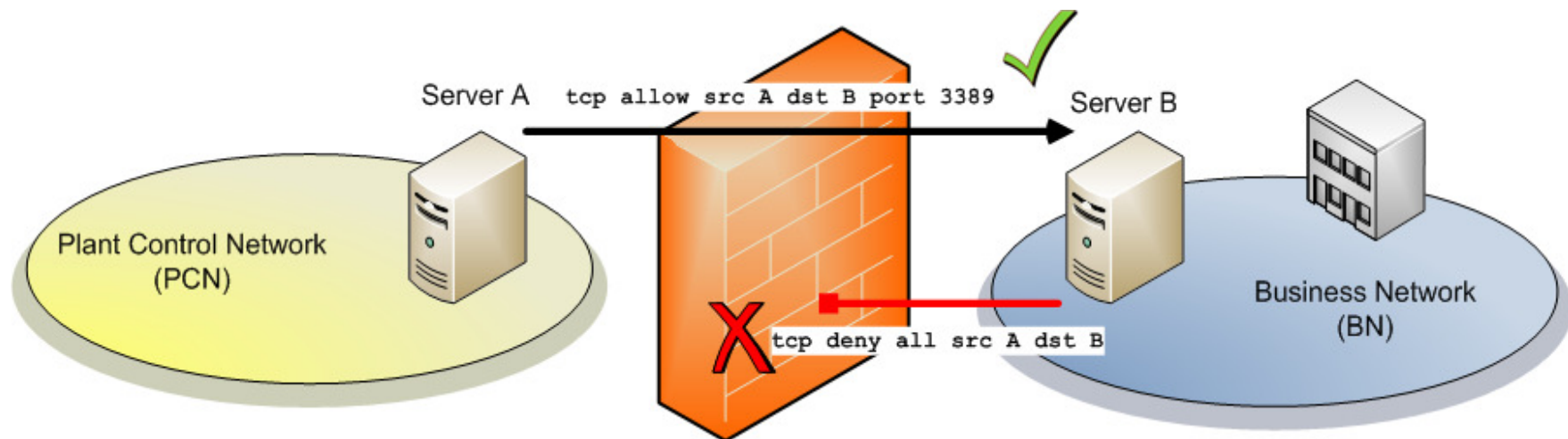
## Both Scenarios:

- Zero-day exploits in desktop software (e.g. browsers, operating system, browser plugin)
- Anti-virus/anti-malware measures will not detect if no signature available
- IDS/IPS will not detect if no signature available or if connection is encrypted
- Payload deploys rootkit or Remote Access Toolkit (RAT)
- Payload initiates outbound connection over SSL/TLS or other encrypted protocol to bypass IDS/IPS/firewall inspection measures
- Attacker now has full control over employee's system and can attack local servers

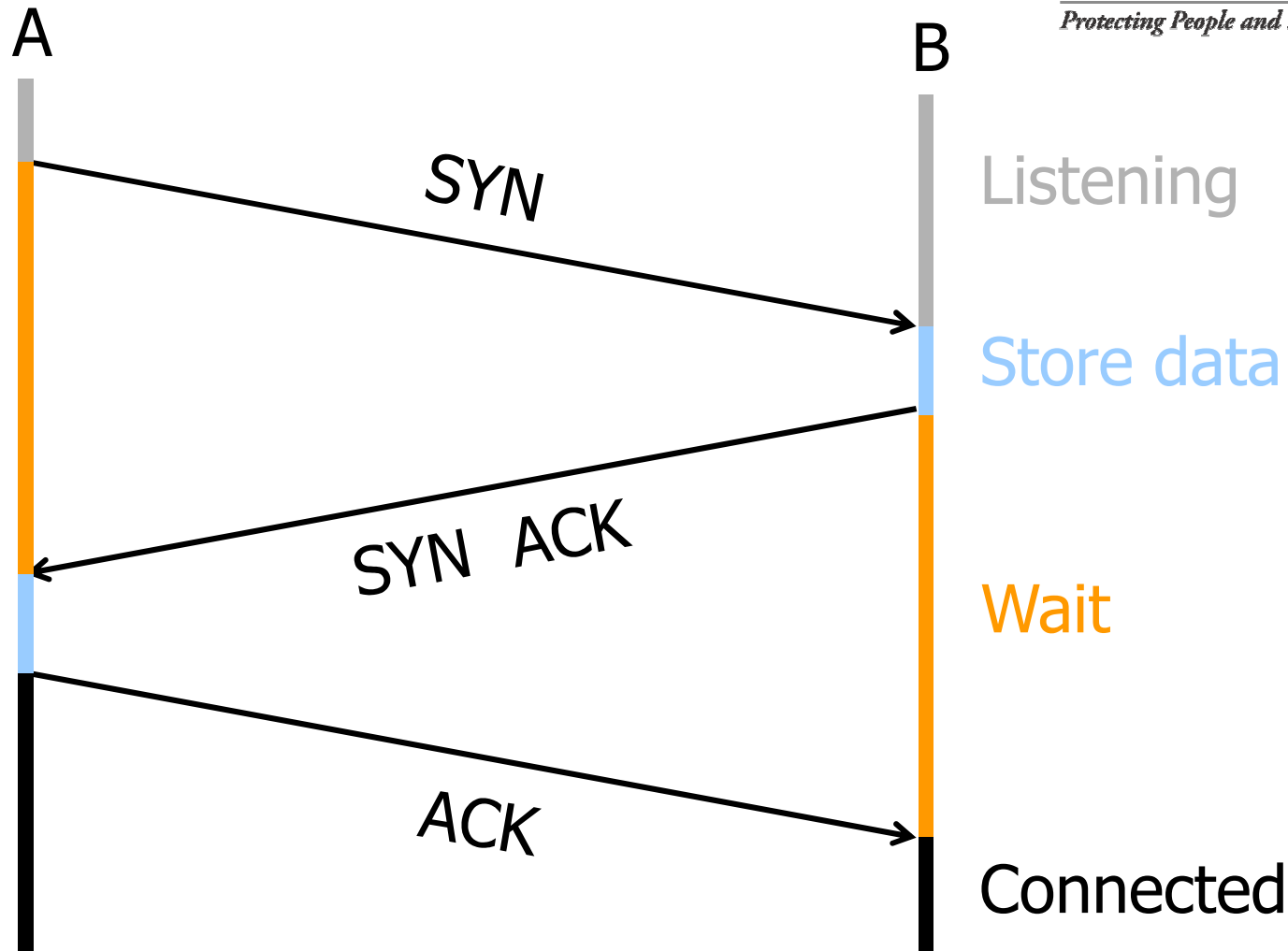
# Attack Stage 2 – Impersonation Attack

## Scenario:

- No connections are allowed thru firewall from PCN to BN
- Firewall is configured as “one way”
- Server A, behind the firewall, sends a requests for data to Server B
- Server B cannot talk to Server A



# TCP “Handshake”

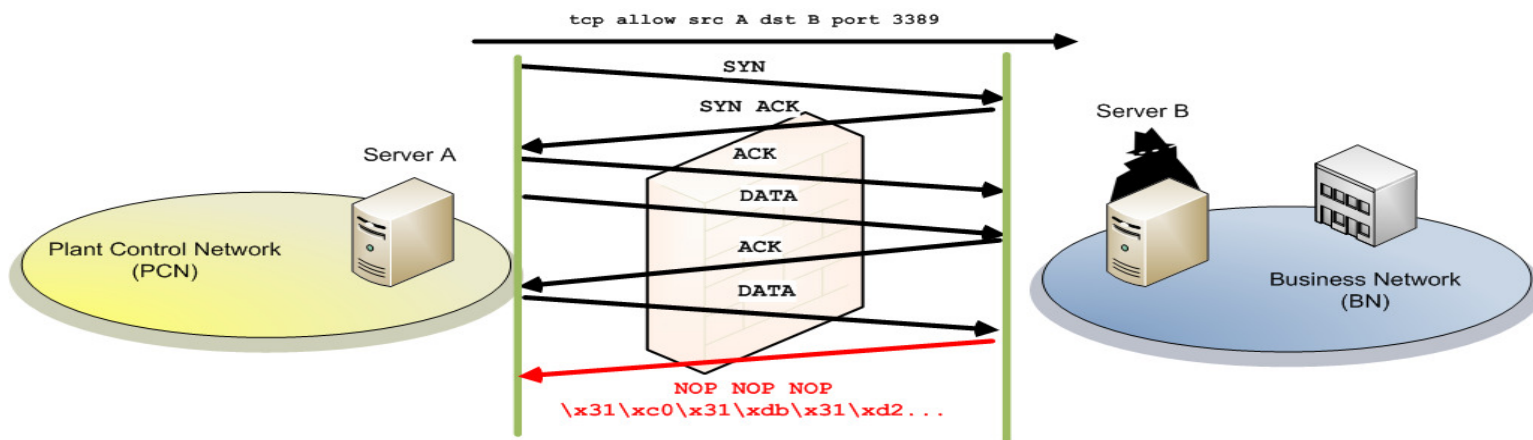


Once established, all TCP connections are bi-directional. Attacks can flow back to clients!

# Attack Stage 2

## Buffer Overflow

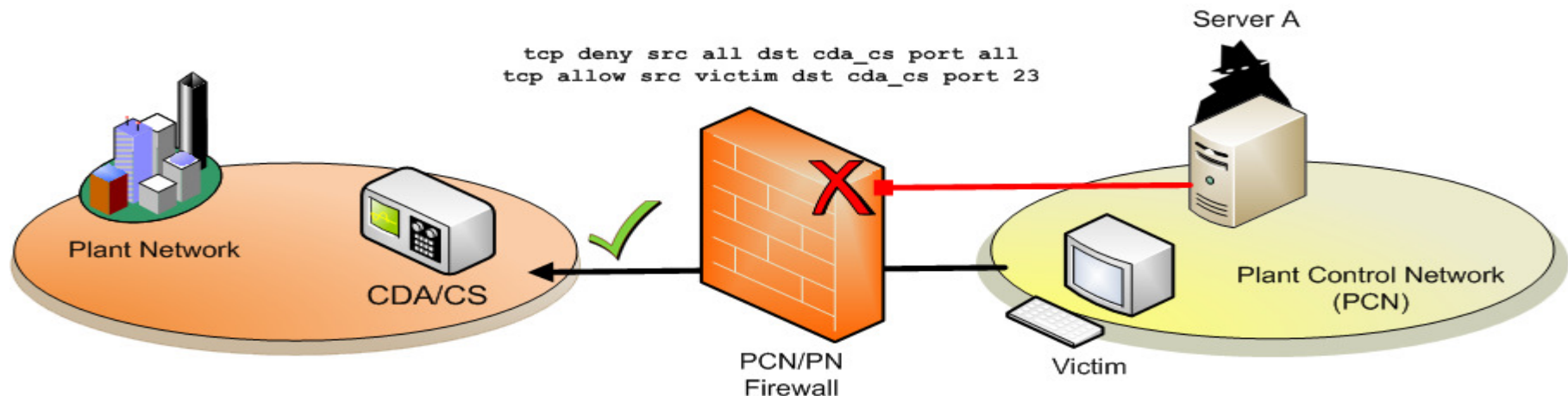
- A buffer overflow occurs when attacker sends data that cannot be adequately handled by the victim program
  - Unexpected value
  - Value out-of-bounds
  - Memory violation
- Attack packet contains executable instructions to request victim open a shell prompt
- The original session has **not** terminated



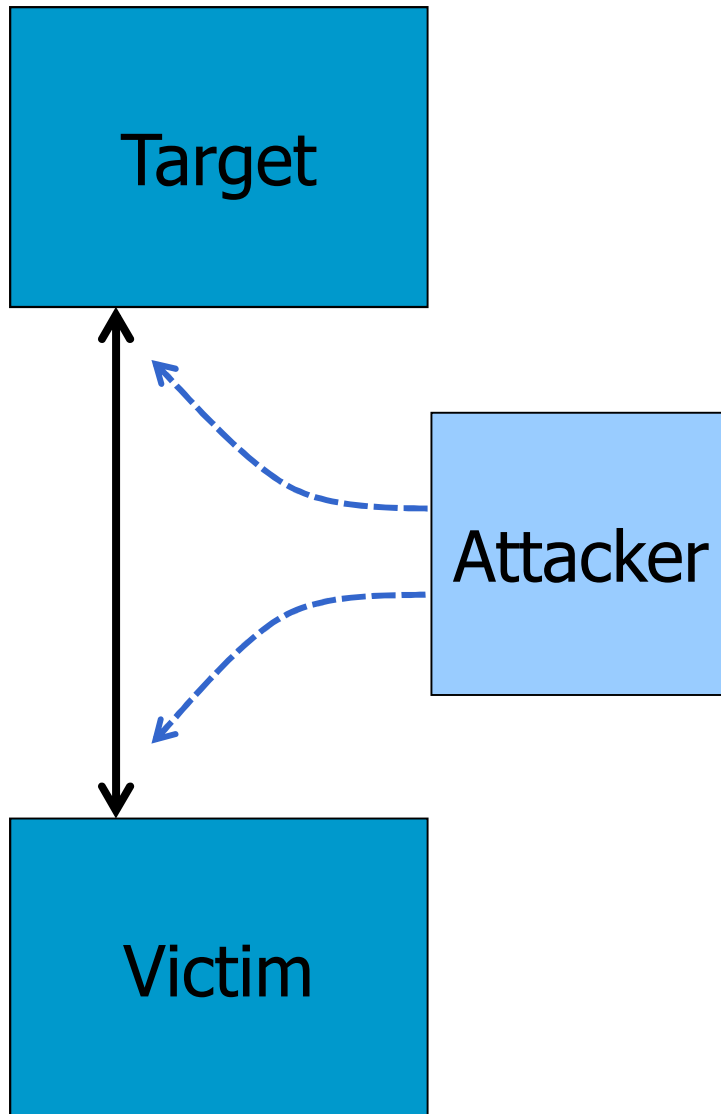
# Attack Stage 3 – Session Hijack

## Scenario:

- Victim is logged into CDA/CS, through the firewall
- Telnet connection is allowed from Victim to ICS
- No other hosts are allowed to connect thru firewall to ICS
- Telnet Connection is authenticated



# Blind TCP Session Hijacking

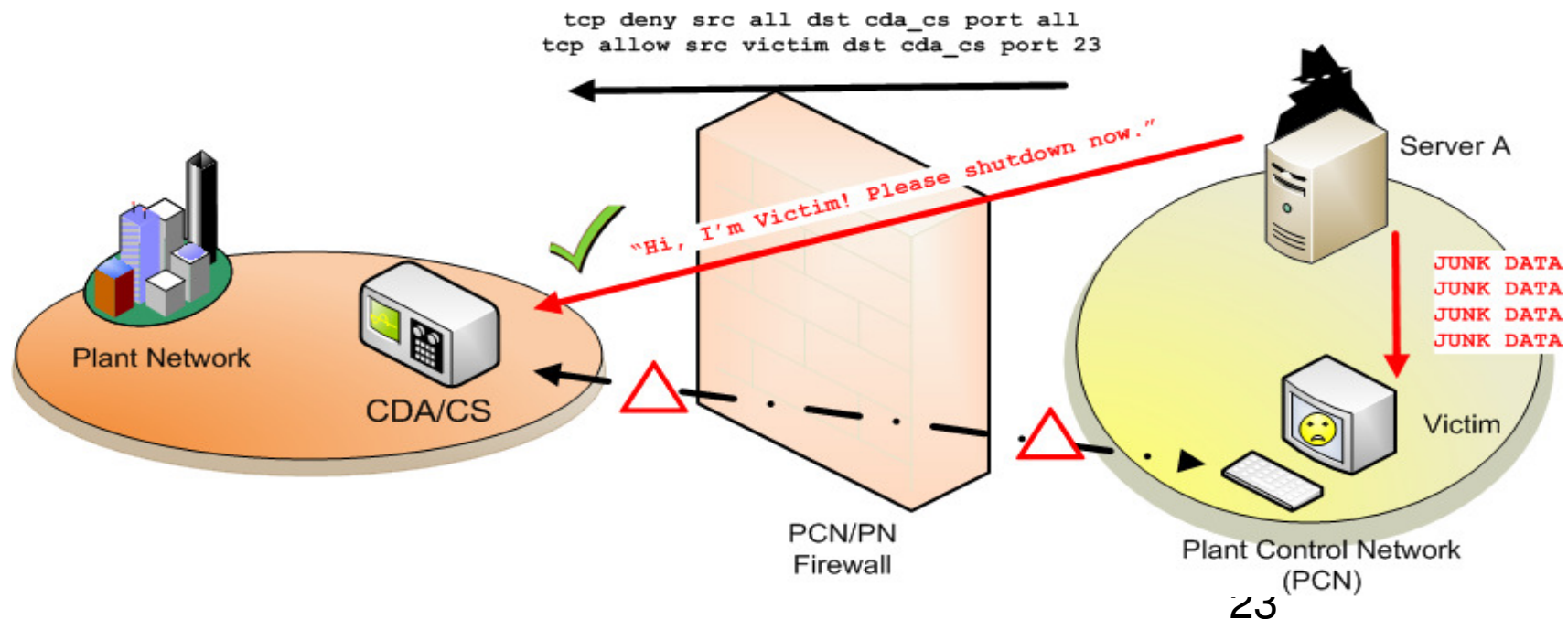


- Victim, target trusted authenticated connection
  - Packets will have predictable sequence numbers
- Attacker impersonates victim to target
  - Opens connection to target to get initial seq number
  - Fills victim's receive queue
  - Sends packets to target that resemble victim's transmission
  - Attacker cannot receive, but may execute commands on target



# Attack Stage 3 – Session Hijack

- Attacker listens to unencrypted session
- Attacker uses probes to determine sequence numbers
- Attacker sends spoofed identity packets to ICS while performing Denial of Service on Victim
- Attacker sends shutdown command to ICS



# How Easy are These Attacks?



- Numerous RAT/trojan toolkits available on underground market
  - Push-button ease of use
  - Exploits as a Service (EaaS) becoming viable business model<sup>1,2</sup>
- Buffer overflow attack methodologies have been well-known and well-documented for many years
  - “Smashing the Stack for Fun and Profit” by AlephOne, *Phrack* magazine, 1996
- Session hijacking is one of the oldest attack methods on the Internet
  - Kevin Mitnick “man-in-the-middle” attack, 1994

1. Grier, Ballard, Caballero, et. al. (2012). Manufacturing Compromise: The Emergence of Exploit-as-a-Service. 19<sup>th</sup> ACM Conference on Computer and Communications Security. Retrieved from <http://cseweb.ucsd.edu/~voelker/pubs/eaas-ccs12.pdf>
2. Asprey, D. (2011). New type of cloud emerges: Exploits as a Service (EaaS). TrendMicro Security. Retrieved from <http://cloud.trendmicro.com/new-type-of-cloud-emerges-exploits-as-a-service-eaas/>

# How Easy are These Attacks?

- Free, easily available hacking tools and toolkits can perform some or all firewall bypass attack types:
  - Metasploit Framework
  - Cain and Abel
  - Firesheep
  - LOIC
  - Evader
  - Backtrack Live CD
  - Nmap
  - Ettercap

# Firewall Limitations



- Firewall technology is not one way (non-deterministic, not application-fluent)
- Firewalls can be bypassed in many ways
- Firewalls have their own vulnerabilities
- Effective Security Programs must do the following:
  - Prevent
  - Detect
  - Delay
  - Deny
  - Deter
  - Respond
  - Recover
- Firewalls cannot do all of these things alone